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Cover Design by Stephen C. Fischer

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# A General Game Playing Program

- BY HERBERT L. DERSHEM

Programmers have used the look-ahead strategy to develop competitive game playing programs for games like checkers and chess. A general form of this look-ahead algorithm can be described in terms of a recursive procedure implemented in BASIC for specific games. If your BASIC processor accepts recursive subroutine calls, then you can use this algorithm to play any suitable game by programming three additional subroutines that describe the game. (For more information on recursive programming, see "Recursive Programming in BASIC", April PC.)

Consider a game with two players called "computer" and "opponent". At any given point in the game, two descriptors describe the situation: the game status (GS), often the status of the game board; and the player to move next (PM), either "computer" or "opponent". Each GS, PM pair results either in a completed game with a winner, or in a draw, or in a set of legal moves for PM. Each legal move maps the GS, PM pair. Let's consider the case where the players alternate moves, making the new PM generated by a move always different from the previous PM.

Now we're ready to recursively state the look-ahead algorithm which, given a GS,PM pair, evaluates all the legal moves available to player PM and determines the optimal one.

Algorithm Evaluate to find the best move BM for player PM from game status GS with evaluation of E. Evaluate (GS,PM,E,BM)

1. If(GS,PM) is directly evaluatable, evaluate it and place result in E; return.

2. Generate  $MV_1$ . $MV_2$ .... $MV_n$ , the set of all legal moves from (GS,PM), and  $GS_1$ , $GS_2$ ,..., $GS_n$ , the corresponding set of game statuses after the legal moves are applied to GS.

3. If PM = computer, call Evaluate  $(GS_i, opponent, E_i, BM_i)$  for i = 1, 2, ..., n; for  $E_k$ , the largest of  $E_1, E_2, ..., E_n$ , set  $E = E_k$ ,  $BM = MV_k$ ; return.

4. If PM = opponent, call Evaluate (GS<sub>i</sub>, computer, E<sub>i</sub>, BM<sub>i</sub>) for i = 1, 2,..., n; for E<sub>k</sub>, the smallest of E<sub>1</sub>, E<sub>2</sub>..., E<sub>n</sub>, set E = E<sub>k</sub>, BM = MV<sub>k</sub>; return.

Evaluation of a game status is always from the computer's point of view. The larger the evaluation, the better the status is for the computer. Therefore, the principle behind this algorithm is that the computer always chooses from the legal moves that move resulting in a game status with largest evaluation. On the other hand, the opponent always chooses the move with the smallest evaluation, since that move is the least desirable for the computer.

How does the computer determine whether a move is directly evaluatable? If a game status is terminal, there are no further moves. Or sometimes the computer stops when a certain number of levels of moves have been examined. For example, a 3-level look-ahead will examine all of the computer's legal responses. As you can see, the number of moves that must be examined grows rapidly as the level of the search infinal level (level 3 in the example above), you must implement some heuristic procedure to evaluate the GS, PM pair. The ability of this procedure, the static evaluation function, to ac-

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curately evaluate the game's status greatly affects how well the computer will compete. There's a trade-off between the depth of look-ahead and the validity of the static evaluation function. If the static evaluation function is perfect, the computer can use it to evaluate all its alternatives directly and not look ahead at all. On the other hand, if the computer can look ahead clear to the end of the game, examining all of the alternatives, it has no need for a static evaluation function since the perfect evaluation function is the game result: win, lose or draw. In practice we find ourselves somewhere between those two extremes.

For the general BASIC version for this algorithm, see Listing 1. Two additions to the algorithm have been made to speed up the search. Both halt the process when it's obvious no more searching is needed.

Suppose the search is at a level generating the computer's responses. If, at the preceding level, the opponent's best move evaluates to 4 and so far the computer's best move at this level evaluates to 5, why continue the search at this level? The opponent will never choose the current move under consideration because it will evaluate to no smaller than 5 which is already 1 worse than the best move the opponent has examined so far. This condition is tested in line 2100 of the program in Listing 1. In tree searching this process, called alphabeta pruning, usually saves search time.

Additional savings can result from statement 2130 where, as soon as a player has found a sure winner for himself, he stops searching.

Now let's look at two implementations of the algorithm in Listing 1. The first, found in Listing 2, is the familiar game of tic-tac-toe. The implementation requires the addition of three subroutines to the general game status evaluator at 2000. These are 1000, a move generator; 3000, a static move evaluation function; and 4000, a gameover tester. But the choices shown here are examples: try designing your own improved versions of these subroutines.

The particular implementation here uses a maximum search depth of 10 levels. For tic-tac-toe, this level implies all searches will be terminated by the end of the game since the longest possible game is 9 moves. The static evaluation function returns 100 if the position is a win for the computer,-100 if it's a win for the opponent, and 0 if it's a draw.

Subroutine 2000 has been modified slightly from that shown in Listing 1 to accommodate the presence of only one subscripted variable in Radio Shack Level 1 BASIC, the system on which this program was implemented.

The ancient game of Kalah, our second game, is played on a board with six small pits on either side and large pits at each end. The game begins with 3 markers in each of the small pits as shown in Figure 1.



#### Figure 1 Initial position of Kalah board

The players alternate moves according to the following rules:

1. A player moves by choosing a pit on his side of the board and distributing the markers contained in that pit into other pits counterclockwise around the board beginning with the counterclockwise neighbor of the emptied pit. He places one marker in each pit and Kalah in turn until all markers removed are distributed. *Example:* If the opponent began play from the initial board shown above by emptying the fifth pit from the left on his side, after his move the board would look like Figure 2.





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2. If the last marker distributed by a player lands in that player's Kalah, the player must empty another pit on his side. This move is called a continuation. The continuation might have another continuation, and so on. *Example:* If the computer's response to the above move was to empty the third pot from the left, it would receive a continuation. This move and its continuation are illustrated in Figure 3.

3. If the last marker distributed on a player's move lands in an empty pit on the player's side of the board, and if some markers are in his opponents pit directly opposite this pit, then the last marker distributed and all the markers in the opposite pit are placed in the Kalah of the player making the move. This move is called a capture. Example: If the opponent now empties the lands in the empty pit and captures the

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computer's four markers on the opposite side (Figure 4).

The winner has the most markers in his Kalah at the end of the game. When a player has no more markers in his pits and it's his turn to move, the game ends. At that point the opponent places all the markers in his pits into his Kalah, and the winner is determined.

Listing 3 shows the application of the game playing algorithm to this game. The continuation complicates matters by requiring two locations to store a move as well as a special coding scheme for continuation moves.



Now that you've seen these examples, you can implement this algorithm for other games. You might want to improve the computer's performance on these games by providing better static evaluation functions or increasing the maximum depth of search. You must proceed with caution, however. Look-ahead algorithms can consume lots of computer time. So be prepared to wait for the computer's moves.

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# Listing 1-Game-Playing Algorithm

1000	1.			
1500	REM	GENERAL GAME PLAYING PROGRAM. THIS SUBROUTINE	1000	
1510	REM	WHICH IS CALLED BY 1990, WILL ACCEPT & GAMP STATUS STOPPD	1000	
1520	REM	IN A(1) THRU A(S) AND RETURN IN S(1) THE BEST MOVE STORED	1000	RI
1530	REM	AND IN E THE EVALUATION OF THAT NOVE THE DUST HOVE FOUND	1900	R
1540	REM	PARAMETERS:	1910	RI
1550	REM	S IS THE NUMBER OF LOCATIONS NEEDED TO STORE THE		DI
1560	REM	M IS THE MAXIMUM DEPTH OF SPARCE	1920	RI
1570	REM	N IS & VALUE WHICH IS INDOSCIDLE CODE FOR A MAN	1930	RI
1575	REM	REPRESENTS A NULL MOVE	1989	RI
1580	REM	WISA VALUE SUCH THAT ANY CARE AT THE	1990	L=
1590	REM	SEW TS A STA FOR THAT ANY GARE STATUS WHICH EVALUATES	1999	RI
1600	REM		2000	L
1610	REM		2009	RE
1620	REM	VARTARI PC.	2010	GC
1630	DEM	T TO TUP TOWN PARAMETER	2018	RS
1640	DPH	I IN THE LEVEL INDICATOR FOR THE CURRENT LEVEL OF SEARCH.	2019	DP
1650	DEM	TADICATES PLAYER WHO IS NOVING: 1=COMPUTER, -1=OPPONENT.	2020	TE
1660	DPW	Q IS THE STACK POINTER. IT INDICATES THE POSITION IN THE	2020	TL
1670	DEM	STACK DIMENSIONED VARIABLE WHERE THE CURRENT GAME STATUS	20 30	
1690	DTH	DESCRIPTION BEGINS.	2040	
1600	8 50	A(L) IS THE CORRENT MOVE BEING EXAMINED AT LEVEL L	2049	RE
1700	REM	S(L) IS THE BEST MOVE EVALUATED SO FAR AT LEVEL L.	2050	
1710	REE	B(L) IS THE EVALUATION OF THE BEST MOVE SO FAR AT LEVEL I	2059	RE
1720	RED	E IS THE VARIABLE IN WHICH THE EVALUATION OF THE BEST MOVE	2060	GO
1720	REA	IS RETURNED.	2009	N E
1700	R5R		2070	Ir
1740	REM	SUBROUTINES:	2079	RE
1750	REM	1000 GENERATES FROM MOVE M(L). THE NEXT MOVE IN A SECURACE	2080	GO
1/60	REM	OF ALLOWABLE MOVES FROM THE GAME STATUS STORED AT DOSTUTION	2086	RE
1770	REM	Q IN THE STACK. THE MOVE IS STORED IN MILL AND THE NEW	2087	RE
1780	REM	GAME STATUS IS PLACED IN THE STACK BEGINNING AT DOST TON	2088	RE
1790	REM	Q+S. THE FIRST MOVE IN THE SECTENCE IS GENERATED UNTER	2089	RE
1800	REM	M(L)=N, THE NULL MOVE, WHEN THE SUBBOUTINE TO CALL BO	2090	IF
1810	REM	IF M(L) IS THE LAST MOVE IN THE SPONDER THE ALLED.	2100	
1815	REM	IS RETURNED.	2108	RE
1820	REM	3000 EVALUATES THE GAME STATUS STORED PROTECTION	2109	RE
1830	REM	O OF THE STACK USING A STATIC PULLINE AND AT POSITION	2110	IF
1840	REM	VALUE IS STORED IN P	2120	1
1850	REM	4000 TESTS THE GAMP STATIC STORES PROVIDENT	2129	RE
1860	REM	OF THE STACK IF IT IS & COMP DEGINNING AT POSITION Q	2130	IF
1870	REM	TE NO MORE MOVES ARE AND AGE ENDING POSITION, THAT IS,	2149	REI
19110	Contraction of the	TO HOME HOTES ARE POSSIBLE, O IS RETURNED AS 1.	3150	(Britten)

- OTHERWISE, O IS RETURNED AS ZERO. IN THIS SUBROUTINE IS WRITTEN IN RADIO SHACK LEVEL I BASIC IN EXCEPT FOR THE USE OF EXTRA IMENSIONED VARIABLES M.S. AND L. IN THESE HAVE BEEN USED FOR CLARITY. M INITIALIZE L AND Z ON THE FIRST CALL. M 0: Z=-1 UPDATE L.Q. AND Z FOR THE NEXT LEVEL OF SEARCH. +1: Q=S\*(L-1) +1: Z=-Z TEST IP GAME IS OVER. SUB 4000 THE HOUSE IS TO THE MAXIMUM OF GAME IS OVER, EVALUATE IN USING STATIC EVALUATION FUNCTION AND RETURN. (L <= 1) + (0=0) goto 2050 goto 2150 NINITIALIZE FOR BEST POSSIBLE MOVE SEARCH. L)=N: S(L)=N: B(L)=-Z\*W M GENERATE NEXT MOVE M GENERATE REAT NOVE SUB 1000 M IF NO MORE MOVES, SET E AND RETURN. M(L) = N THEN E=B(L): GOTO 2150 M EVALUATE THIS MOVE, M(L), BY A RECURSIVE CALL. SUB 2000 M IF THE BEST MOVE AT THIS LEVEL IS ALREADY RETTER FOR Z M THAN THE BEST MOVE FROM PRECEDING LEVEL WAS FOR -Z, THEN M THAN THE BEST MOVE PROM DECODEN BY -Z ANYWAY, SO RETURN M WITHOUT EVALUATING THE OTHER MOVES AT THIS LEVEL. L=1 GOTO 2110 IF 2\*E>=2\*B(L-1) THEN B(L)=E: GOTO 2150 M IF THIS IS THE FIRST RESPONSE THED OR IT IS BETTER H THAN THE BEST SO PAR, RECORD IT AS BEST SO PAR. (5(L)=E: S(L)=H(L)
  - (5(L)(5))\*((2\*E(=2\*E(L)) GOTO 2060 B(L)=E: S(L)=H(L) M IF THIS RESPONSE WINS, THERE IS NO NEED TO SEARCH MORE. 2\*B(L) < W GOTO 2060 (M ADJUST L.Q. AND Z AND RETURN. -L-1: Q=Q-S: Z=-Z: RETURN
  - 150 L=



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# Listing 2 – Tic-Tac-Toe

954 REM TIC-TAC-TOE MOVE EVALUATOR IN RADIO SHACK LEVEL I BASIC.	2030 GO
956 REM THE BOARD POSITION IS SORED IN A(O) THEU A(0+8) AS	2040 GO
958 REM X X	2050 A (Q+
959 REM A(Q) X A(0+1) X A(0+2)	2060 GOSU
960 REM XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX	2070 IF A
961 REM A (0+3) X A (0+4) X A (0+5)	2080 GOSU
962 REM XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX	2090 IF L
963 REM A(0+6) X A(0+7) X A(0+8)	2100 IF Z
964 REM X X	2110 IF (
965 REM	2120 A (Q+
968 REM & POSITION UNOCCUPIED CONTAINS & O.	2130 IF Z
972 REM & FOSITION OCCUPIED BY THE COMPUTER'S MARK CONTAINS & 1	2150 L=L-
974 REM & FOSITION OCCUPTED BY THE OPPONENT'S MARK CONTAINS & 4	2986 REM
975 REM M(L), S(L), AND B(L) FROM THE GENERAL ALGORITHM (SEE FIG	1) 2988 REM
976 REM ARE STORED IN A (0+9) A (0+10) AND A (0+11) DESPECTIVELY	2990 REM
978 REN VALUES OF PARAMETERS:	2992 REM
980 REM S=12	2994 REM
982 REM M=10 (SEARCHES UNTIL COMPLETION OF GAME)	2996 REM
984 REM N=0	2998 REM
986 REM W= 100	3000 GOSU
988 REM	3010 IF A
990 REM	3020 IF V
992 REM SUBROUTINE 1000 STORES IN A (0+9) THE NEXT MOVE FOR BOARD	3030 E=-0
994 REM POSITION A (0) - A (0+8) FROM PREVIOUS MOVE A (0+9). IF A (0+9)	=0. 3988 REM
996 REM FIRST MOVE IS RETURNED. IF THERE ARE NO MORE MOVES, A (0+9)	) 3990 REM
998 REM IS RETURNE AS ZERO. NEW BOARD IS STORED IN A (0+S) -A (0+S+8	3992 REM
$1000 \ A (0+9) = A (0+9) + 1$	3994 REM
1010 IF A(0+9)>9 THEN A(0+9)=0: RETURN	3996 REM
1019 REM IF POSITION IS OCCUPIED, TRY THE NEXT ONE.	3998 REM
1020 IF A (0+A (0+9) -1) <>0 GOTO 1000	3999 REM
1030 FOR I=0 TO 8	4000 REST
1040  A(Q+S+I) = A(Q+I)	4010 FOR
1050 NEXT I	4020 RE
1059 REM RECORD THE MOVE.	4030 T=
$1060 \lambda (Q+S+\lambda (Q+9)-1) = (Z=1)+4*(Z=-1)$	4040 IF
107C RETURN	4050 IF
1986 REM	4060 IF
1988 REM TIC-TAC-TOE VERSION OF GENERAL EVALUATION ALGORITHM	4070 NEXT
1990 L=0: Z=-1: S=12: M=10: N=0: W=100	4080 O= (V
2000 L=L+1:Q=5*(L-1)+1: Z=-Z	4090 REIU
2010 GOSUB 4000	4099 REM
2020 IF (L<=M)* (0=0) GOTO 2050	4100 DATA

2030 GOSUB 3000 2040 GOTO 2150 2055 A(Q+9)=N: A(Q+10)=N: A(Q+11)=-Z\*M 2060 GOSUB 1000 2070 IF A(Q+9)=N THEN E=A(Q+11) GOTO 2150 2080 GOSUB 2000 2090 IF L=1 GOTO 2110 2100 IF Z#D>Z=X (Q-5+11) THEN A(Q+11)=E: GOTO 2150 2110 IF (A(Q+10)<=N)\*(Z\*E<=Z\*A(Q+11)) GOTO 2060 2120 A(Q+11)=E: A(Q+10)=A(Q+9) 2130 IF Z\*A(Q+11)(X\*G GOTO 2060 2150 L=L-1: Q=Q-S: Z=-Z: RETURN 2968 REM STATIC MOVE EVALUATOR FOR TIC-TAC-TOE. 2990 REM E=100 IF WINNING POSITION FOR THE COMPUTER. 2998 REM E=-100 IF WINNING POSITION FOR THE COMPUTER. 2998 REM E=-0.5 IF NOT A GAME-OVER POSITION. 3000 GOSUB 4000 3010 IF ABS(Y)=100 THEN E=V: RETURN 3020 IF V=8 THEN Z=0: RETURN 3030 E=-0.5: RETURN 3030 E=-0.5: RETURN 3040 HEM 0=1, V=100 IF WINNING POSITION FOR THE OPDONENT. 3999 REM 0=1, V=100 IF WINNING POSITION FOR COMPUTER. 3990 REM SUBROUTINE EXAMINES A(Q)-A(Q+8) AND RETURNS: 3994 REM 0=1, V=100 IF WINNING POSITION FOR COMPUTER. 3995 REM 0=1, V=0 IF DRAW POSITION. 3999 REM 0=1, V=0 IF NOT A GAME ENDING POSITION. 3999 REM 0=1, V=0 IF NOT A GAME ENDING POSITION. 3999 REM 0=1, V=0 IF NOT A GAME ENDING POSITION. 3999 REM 0=0, V.C8 IF NOT A GAME ENDING POSITION. 3999 REM 0=0, V.C8 IF NOT A GAME ENDING POSITION. 3999 REM 0=0, V.C8 IF NOT A GAME ENDING POSITION. 3999 REM 0=0, V.C8 IF NOT A GAME ENDING POSITION. 3999 REM 0=0, V.C8 IF NOT A GAME ENDING POSITION. 3999 REM 0=0, V.C8 IF NOT A GAME ENDING POSITION. 3999 REM 0=0, V.C8 IF NOT A GAME ENDING POSITION. 3990 REM 0=1, V=0 IF NO IF WINNING POSITION. 3990 REM 0=1, V=0 IF NO IF NEINING POSITION. 3990 REM 0=1, V=0 IF NO IF NEINING POSITION. 3990 REM 0=1, V=0 IF NO IF WINNING POSITION. 3990 REM 0=1, V=0 IF NO IF



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# Listing 3-Kalah

938 REM KALAH GAME EVALUATOR IN RADIO SHACK LEVEL I BASIC. 940 REM THE BOARD IS STORED IN A(Q) THRU A(Q+13) AS 942 REM 942 HER 944 REM 946 REM A(Q+12) A(Q+11) A(Q+10) A(Q+9) A(Q+8) A(Q+7)948 REM A(Q+13) A(Q+11) A(Q+10) A(Q+9) A(Q+8) A(Q+7)950 REM 950 REM 9510 REM 952 REM 954 REM 954 REM 956 REM A SIMPLE MOVE IS REPRESENTED BY AN INTEGER 0-5 WITH 0 956 REM A SIMPLE MOVE IS REPRESENTED BY AN INTEGER 0-5 WITH 0 958 REM REPRESENTING THE PIT FARTHEST FROM THE PLAYER'S KALAH 960 REM AND 5 THE PIT NEAREST THE PLAYER'S KALAH. 962 REM 964 REM AS FOLLOWS: 965 REM (MOVED1\*6(I-1) + (MOVED1\*6(I-2)) 965 REM (MOVEI) = 6(T-1) + (MOVE2) + 6(T-2) + ... + (MOVEI) + 60.965 REM THE LCCATION FOLLOWING THIS CONTAINS 6(T-1) TO INDICATE 968 REM THE NUMBER OF CONTINUATIONS. 970 REM M(L) IS STORED IN A (Q+14) , A (Q+15) 974 REM S(L) IS STORED IN A (Q+16) , A (Q+17) 974 REM S(L) IS STORED IN A (Q+18) 976 REM VALUE OF PARAMETERS: 978 REM S=19 980 REM 982 REM M IS UNDER EXTERNAL PROGRAM CONTROL 982 REM N=-1 984 REM N=100 986 REM 98100 986 REM SUBROUTINE 1000 STORES IN A (Q+14), A (Q+15)THE NEXT MOVE FOR 990 REM BOARD POSITION A (Q) -A (Q+13) FROM PREVIOUS MOVE A (Q+14), 992 REM A (Q+15). IF A (Q+15)=-1, THE FIRST MOVE IS RETURNED, 994 REM IF THERE ARE NO MORE MOVES, A (Q+14) IS RETURNED AS -1. 994 REM IF THERE ARE NO MORE MOVES, A (Q+14) IS RETURNED AS -1. 996 REM THE RESULTING BOARD POSITION IS STORED IN A (Q+S) -A (Q+S+13). 994 REM IF THERE ARE NO MORE MOVES, A(Q+14) IS RETURNED AS -1. 996 REM THE RESULTING BOARD POSITION IS STORED IN A(Q+S) - A(Q+S+13). 997 REM 998 REM INCREMENT MOVE AND STORE IN T AND R. 1000 A(Q+14) = A(Q+14) + 1: T = A(Q+14): R = A(Q+15)1000 POR I=0 TO 13: A(Q+S+1) = A(Q+1): NEXT I1010 POR I=0 TO 13: A(Q+S+1) = A(Q+1): NEXT I1019 REM IF MOVE AT ONE CONTINUATION EXHAUSTED, COME BACK A LEVEL. 1020 IF (INT(T/5)\*5=T)\*(R<>1) THEN T=T/6: R=R/6: GOTO 1020 1024 REM TEST FOO LAST HOVE. 1025 IF (T=6)\*(R=1) THEN A(Q+14) = -1: RETURN1030 U=T: V=R 1030 REM PULL OUT SIMPLE MOVE. 1049 REM P IS THE PIT POSITION ON BOARD OF MOVE. 1050 P=7\*(Z=1)+X 1055 REM IF PIT IS EMPTY, GO GET ANOTHER HOVE. 1060 IF A(Q+P+S) = 0 THER T=T+1: GOTO 1010 1069 REM MAKE THE MOVE. 1070 D=Q+S: FOR I=P+1 TO P+A(D+P) 1080 A(D+3) = A(D+2) = 01110 NEAT I 1110 I=A(D+P): A(D+P) = 01119 REM IF A NEW CONTINUATION IF SOUND, GO FORWARD A LEVEL. 1200 IF (J=6+7\*(Z=1)) \* (Y=0) THEN T=T\*6: R=R\*6: U=0: Y=1 1050 A(0+3) = A(0+3) + 1 100 B(T) = A(0+3) = A(0+3) + 1 1100 E(T) = A(0+2) = A(0+3) + 1 1110 E(T) = A(0+2) = A(0+3) + A(0+3) = A(0+3) + A(0+3) 2982 REM 2984 REM STATIC MOVE EVALUATOR FOR KALAH. 2966 REM E=100 IF WINNING POSITION FOR THE COMPUTER. 2986 REM E=-100 IF WINNING POSITION FOR THE COMPUTER. 2987 REM ELSE: 2988 REM E= (CONTENT OF COMP'S KALAH) - (CONTENTS OF OPP'S KALAH))\* 2980 REM (1 + 1/(19 - MAX CONTENTS OF A KALAH)) 2990 REM (1 + 1/(19 - MAX CONTENTS OF A K) 2996 REM WINNING POSITION? 3000 GOSUB 4000 3010 E=A4(Q+13) - A(Q+6) 3020 IF 0=1 THEN E=100\*((E>0) - (E<0)): RETURN 3030 F=A(Q+13) + (E>0) + A(Q+6)\*(E<0) 3040 E=E\*(1+1/(19-F)) 3050 RETURN 3992 REM GANE-OVER TESTER POR KALAH. 3996 REM 0=1 IF GAME IS OVER. 4000 REM TEST FOR A WINNER. 4000 0=0 4000 REM TEST FOR A WINNER. 4010 IP (A(Q+13)>18)+(A(Q+6)>18) THEN 0=1: RETURN 4019 REM TEST FOR NOVER'S FITS ALL EMPTY. 4020 J=7+(Z=1) 4030 FOR I=J TO J+5 4040 IF A(Q+1)<>0 THEN RETURN 4050 NEXT I 4059 REM THEY ARE, SO EMPTY OTHER'S FITS INTO KALAH. 4060 FOR I=7-J TO 12-J 4070 A(Q+13-J)=A(Q+13-J)+A(Q+I): A(Q+I)=0 4080 NEXT I 4080 NEXT I 4090 0=1: RETURN



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